Satisficing Approach to Human-in-the-Loop Safeguarded Control

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Abstract— In this paper, control Lyapunov functions are used to define static and dynamic safe regions for a system. Based on a control Lyapunov function, a "satisficing input set" is defined as all satisficing controls for this CLF according to attractive behaviors and repulsive behaviors. If the system is within a specified safe region, the human input will be used as the control input to the system. If the system is outside the specified safe region, the human input will be snapped to the closest control element in the satisficing input set. Behavior based strategies are applied to achieve smooth transition from human input to snapped control input so as to guarantee maximum flexibility for humans as well as system stability and minimum base-line performance.

I. INTRODUCTION

Human-machine interaction as an attractive research area has been studied extensively in the literature with both military and civilian applications. Following the taxonomy in [1], human-machine interaction can be categorized as teleoperation, shared control, traded control, and supervisory control.

The main purpose of this paper is to propose a safeguarded control scheme with a human in the loop which not only allows maximum flexibility for humans but also guarantees system stability. This paper falls into the category of safeguarded shared control. While most current approaches to shared control are based on heuristic principles that are not easily analyzed, our approach is built on analytically rigorous control theory. This approach rests on several seemingly loosely related theoretical foundations, that is, control Lyapunov functions (CLFs) [2], CLF based satisficing control [3], shared control frameworks [4], [5], and reactive control techniques [6], [7], [8].

In this paper, we will use control Lyapunov functions to define static and dynamic safe regions for a system. Based on a control Lyapunov function, a "satisficing input set" can be defined accordingly, which corresponds to all satisficing controls for this CLF. If the system is within a specified safe region, the human input will be used as the control input to the system even if it is outside the satisficing input set. If the system is outside the specified safe region, the human input will be snapped to the satisficing input set. In this case, the actual control input to the system will be the closest control element in the satisficing input set. Behavior based strategies are applied to guarantee smooth transition from human input to snapped control input. Therefore, this approach gives the human maximum flexibility within the constraints that system stability and minimum base-line performance are guaranteed. This paper applies the concept of CLF to find satisficing controls according to attractive behaviors and repulsive behaviors. In this sense, this paper

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can be thought of as a novel use of potential field like approaches (e.g. [8]) to compute safe controls for a vehicle by allowing humans to control the vehicle as long as safety constraints are not violated. The main contribution of this paper is to introduce a rigorous analytic approach to analyzing safeguarded shared control problems using Lyapunov techniques. Instead of restricting our attention to navigation or obstacle avoidance scenarios, we focus on the stability issue of general nonlinear systems with a human in the loop; for example, vehicles with complicated nonlinear dynamics or unstable modes. In fact, system stability is a critical concern for any control loop with humans involved. This feature also distinguishes our approach from the previous shared control frameworks and reactive control techniques.

II. CLF BASED SATISFICING INPUT SET

In this section, control Lyapunov functions will be used to define satisficing input sets for both time-invariant and time-varying affine nonlinear systems.

Consider a time-varying affine nonlinear system

$$\dot{x} = f(x,t) + g(x,t)u, \tag{1}$$

where $x \in \mathbb{R}^n$, $u \in \mathbb{R}^m$, and $f : \mathbb{R}^n \times \mathbb{R}_+ \to \mathbb{R}^n$ and $g : \mathbb{R}^n \times \mathbb{R}_+ \to \mathbb{R}^{n \times m}$ are locally Lipschitz in x and piecewise continuous in t.

Similarly, a time-invariant affine nonlinear system is given by

$$\dot{x} = f(x) + g(x)u. \tag{2}$$

We consider two competing behaviors including an attractive behavior and a repulsive behavior.¹ In the case of attractive behaviors, satisficing controls correspond to all stabilizing controls with respect to a particular CLF which can drive x to a reference trajectory $x_r(t)$ for system (1) or regulate state x to a certain state x^* for system (2). In the case of repulsive behaviors, satisficing controls correspond to all destabilizing controls with respect to a particular CLF which can drive state x away from a reference trajectory $x_r(t)$ for system (1) or from a certain state x^* for system (2). With regard to these two behaviors, we have the following control Lyapunov function definition for system (1).

Definition 2.1: A continuously differentiable function $V : \mathbb{R}^n \times \mathbb{R}_+ \to \mathbb{R}$ is a control Lyapunov function (CLF)

¹Although terms "attractive behavior" and "repulsive behavior" are used exclusively in robot navigation and motion planning literature for goal seeking and obstacle avoidance, here we employ the same terms but refer to general constraints for the system state. In this paper, attractive behavior corresponds to the case when some nonnegative function of the state is constrained to be below a certain upper bound while repulsive behavior corresponds to the case when some nonnegative function of the state is constrained to be above a certain lower bound.

for system (1) if it is positive definite, decrescent, radially unbounded in x, and satisfies

$$\inf_{u} \left\{ \rho \frac{\partial V}{\partial t} + \rho \frac{\partial V}{\partial x} \left(f(x,t) + g(x,t)u \right) \right\} \le -W(x), \quad (3)$$

 $\forall x \neq 0$ and $\forall t \geq 0$ where $\rho = 1$ or -1 and W(x) is a continuous positive definite function. Function V corresponds to a stabilizing CLF for $\rho = 1$ and a destabilizing CLF for $\rho = -1$.

The satisficing input set for system (1) is defined as

$$F(x) \stackrel{\triangle}{=} \left\{ u : \rho \frac{\partial V}{\partial t} + \rho \frac{\partial V}{\partial x} \left(f(x,t) + g(x,t)u \right) \le 0 \right\},\tag{4}$$

where $\rho = 1$ or -1. The set F(x) corresponds to a stabilizing set for $\rho = 1$, denoted as S(x), and a destabilizing set for $\rho = -1$, denoted as $\bar{S}(x)$. Similarly, the asymptotically satisficing input set for system (1) is defined as

$$F_a(x) \stackrel{\triangle}{=} \left\{ u : \rho \frac{\partial V}{\partial t} + \rho \frac{\partial V}{\partial x} \left(f(x,t) + g(x,t)u \right) \le -W(x) \right\}.$$
(5)

Note that both F(x) and $F_a(x)$ are guaranteed to be nonempty by Definition 2.1.

For system (2), CLF and satisficing input set are defined similarly except that CLF V(x) is not an explicit function of time.

The value of ρ is specified according to design purposes. For the particular example of vehicle navigation, we assume that V = 0 at a certain goal and V qualitatively represents how far a vehicle is from its goal. If we want the vehicle to move to this goal, ρ is specified as 1, that is, controls corresponding to negative \dot{V} are feasible. If we want the vehicle to move away from this goal, ρ is then specified as -1, that is, controls corresponding to positive \dot{V} are feasible.

Remark 2.2: Note that Definition 2.1 can be extended to a weak control Lyapunov function (wCLF) definition by allowing V to be positive semi-definite, which may be appropriate for certain applications where only part of the states are of interest.

In the case of $u \in \mathbb{R}^2$, at each time t, the line defined by $\dot{V} = 0$ separates the control space (a plane in this case) into two halves, where one half plane represents the satisficing input set and the other half represents the unfeasible input set. Of course, if $\frac{\partial V}{\partial x}g = [0,0]$ at a certain time t_1 , the satisficing input set will be the whole plane given the existence of the CLF.

We have the following lemma regarding the connection between the stabilizing set S(x) and the destabilizing set $\bar{S}(x)$.

Lemma 2.1: Given a certain CLF without input constraints, if $\frac{\partial V}{\partial x}g(x) \neq 0$ for all $x \neq 0$, then a function V is a stabilizing CLF (destabilizing CLF) if and only if it is also a destabilizing CLF (stabilizing CLF). That is, both S(x) and $\bar{S}(x)$ are nonempty for all $x \neq 0$. Otherwise, if a function V is a stabilizing CLF (destabilizing CLF), it cannot be a destabilizing CLF (stabilizing CLF).

Proof: If $\frac{\partial V}{\partial x}g(x) \neq 0$ for all $x \neq 0$, there always exists a control $u \in \mathbb{R}^m$ to make \dot{V} negative or positive for all

 $x \neq 0$, that is, both S(x) and $\overline{S}(x)$ are nonempty for all $x \neq 0$.

If $\frac{\partial V}{\partial x}g(x) = 0$ for some $x \neq 0$, a stabilizing CLF implies $\frac{\partial V}{\partial x}g = 0 \Longrightarrow \frac{\partial V}{\partial t} + \frac{\partial V}{\partial x}f < 0$. Thus this CLF cannot be a destabilizing CLF. Similarly, a repulsive CLF cannot be a stabilizing CLF in this case.

To illustrate Lemma 2.1, consider the following system

$$\dot{x}_1 = -x_1 + u, \qquad \dot{x}_2 = -x_2 - u.$$
 (6)

It is straightforward to show that $V = \frac{1}{2}(x_1^2 + x_2^2)$ is an stabilizing CLF for system (6). The derivative of V is given by $\dot{V} = -x_1^2 - x_2^2 + (x_1 - x_2)u$. In the case of $x_1 = x_2 \neq 0$, the stabilizing set $S(x) = I\!R$ while the destabilizing set $\bar{S}(x) = \emptyset$. That is, in general nonempty S(x) does not imply that $\bar{S}(x)$ is nonempty, and vice versa.

Note that both the satisficing input set and asymptotically satisficing input set can be extended easily to the case when there are input constraints for the system if a corresponding constrained CLF is known.

III. SATISFICING APPROACH TO SAFEGUARDED SHARED CONTROL

In this section, we propose a satisficing approach to safeguarded shared control. The basic idea behind this approach is that both system stability and minimum levels of performance are guaranteed in the presence of human inputs.

A. Safeguarded Control for Static Safe Region

Motivated by the properties of CLF, we use it to define safe regions for the shared control problem.

Definition 3.1: A static safe region corresponds to a region of the state space where the system has guaranteed minimum base-line performance.

A simple example of a static safe region for an airplane may correspond to a region of its states (e.g. range of roll angles, velocity, and so forth) where the airplane can fly safely.

Safe regions for system (2) are defined for two different behaviors. In the case of attractive behaviors, a safe region is defined as $R = \{x : V(x) \le c\}$, where c is a positive constant and can be specified based on the design requirement, and V(x) is a stabilizing CLF (wCLF) for system (2). In the case of repulsive behaviors, a safe region is defined as $R = \{x : V(x) > c\}$, where c is defined similarly and V(x)is a destabilizing CLF (wCLF) for system (2). In both cases, we assume that the CLF (wCLF) V(x) = 0 at $x = x^*$, where x^* is the reference state as mentioned in Section II. For example, in a vehicle navigation scenario, x^* can be designed as the goal for the goal seeking behavior or as the center of the hazardous region for the obstacle avoidance behavior. For general cases, x^* can be designed as the state reference to which a system needs to be stabilized or as the state reference of a certain unstable mode for systems with unstable modes to guarantee system stability. If we associate different CLFs (wCLFs) to a system for different design purposes, the overall safe region for the system can be designed correspondingly as an intersection of a group of the two kinds of subregions defined above. We have the following definition for the construction of the overall safe region.

Definition 3.2: Given known CLFs V_i and positive scalars c_i , $i = 1, \dots, K$, the overall static safe region for system (2) can be constructed as $R_s(c) = \left(\bigcap_{i=1}^J R_i(c_i)\right) \bigcap \left(\bigcap_{i=J+1}^K \bar{R}_i(c_i)\right)$, where $c = [c_1, \dots, c_K]^T, R_i, i = 1, \dots, J$, represents the attractive subregion $\{x : V_i(x) \leq c_i\}$, and \bar{R}_i , $i = J + 1, \dots, K$, represents the repulsive subregion $\{x : V_i(x) \geq c_i\}$.

Let $\mathcal{I} = \{1, \dots, J\}$ and $\mathcal{I} = \{J + 1, \dots, K\}$. Accordingly, define the *i*th satisficing input set corresponding to the *i*th safe subregion as

$$U_{i}(x) = \begin{cases} \{u : u \in I\!\!R^{m}\}, & x \in R_{i}(c_{i} - \epsilon_{i}) \\ \{u : \dot{V}_{i}(x) \leq 0\}, & x \notin R_{i}(c_{i} - \epsilon_{i}) \end{cases} \quad i \in \mathcal{I} \\ \bar{U}_{i}(x) = \begin{cases} \{u : u \in I\!\!R^{m}\}, & x \in \bar{R}_{i}(c_{i} + \epsilon_{i}) \\ \{u : \dot{V}_{i}(x) \geq 0\}, & x \notin R_{i}(c_{i} + \epsilon_{i}) \end{cases} \quad i \in \bar{\mathcal{I}}, \end{cases}$$
(7)

where ϵ_i represents the tolerance for safeguarding with $0 < \epsilon_i < c_i, i \in \mathcal{I}$, and $\epsilon_i > 0, i \in \overline{\mathcal{I}}$.

We have the following definition for the construction of the overall satisficing input set.

Definition 3.3: Given the overall safe region defined in Definition 3.2, the overall satisficing input set for system (2) can be constructed as $U = (\bigcap_{i=1}^{J} U_i(x)) \bigcap (\bigcap_{i=J+1}^{K} \overline{U}_i(x))$, where $U_i(x)$ and $\overline{U}_i(x)$ are given by (7).

Note that in the above lemma for the overall satisficing input set, we did not consider the possible input constraints for the system, which can be incorporated conveniently in Eq. (7) by constraining $u \in \mathcal{U}$, where \mathcal{U} is the constraint set for the input. Also, the above overall satisficing input set can be easily extended to the overall asymptotically satisficing input set, denoted as U_a , following Eq. (5).

To illustrate the above safe region and satisficing input set concepts, we consider a navigation scenario when a human drives a vehicle in a room. We want to give the human authority and flexibility to drive the vehicle but modify or override his or her control under hazardous situations (e.g. collision). For illustrative purpose, three CLF level curves are plotted in Figure 1. We use a wedge to represent the vehicle. Let $V_i([x, y])$, i = 1, 2, 3, where (x, y) is the Cartesian position of the vehicle, be the CLF corresponding to goal (x_i^*, y_i^*) respectively. We assume that $V_i = 0$ at each goal point. We also assume that the vehicle should be constrained inside the level curve $V_1 = c_1$ (e.g. wall avoidance) as well as outside level curves $V_2 = c_2$ and $V_3 = c_3$ (e.g. obstacle avoidance). In this case, the safe region is defined as $R = R_1 \bigcap \overline{R}_2 \bigcap \overline{R}_3$, where $R_1 =$ $\{x : V_1(x) \leq c_1\}, \bar{R}_2 = \{x : V_2(x) \geq c_2\}, \text{ and }$ $\bar{R}_3 = \{x : V_3(x) \ge c_3\}$. Accordingly, the overall satisficing

input set is defined as $U = U_1(x) \cap \overline{U}_2(x) \cap \overline{U}_3(x)$, where

$$\begin{split} U_1(x) &= \begin{cases} \{u: u \in I\!\!R^2\}, & x \in R_1 \\ \{u: \dot{V}_1(x) \leq 0\}, & x \notin R_1 \end{cases}, \\ \bar{U}_2(x) &= \begin{cases} \{u: u \in I\!\!R^2\}, & x \in \bar{R}_2 \\ \{u: \dot{V}_2(x) \geq 0\}, & x \notin \bar{R}_2 \end{cases}, \\ \bar{U}_3(x) &= \begin{cases} \{u: u \in I\!\!R^2\}, & x \in \bar{R}_3 \\ \{u: \dot{V}_3(x) \geq 0\}, & x \notin \bar{R}_3 \end{cases}. \end{split}$$

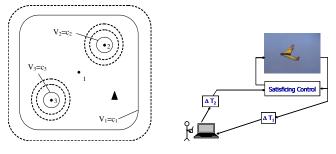


Fig. 1. CLF level curves.

Fig. 2. A UAV equipped with an onboard satisficing controller.

Consider another scenario where a pilot controls an unmanned air vehicle via a ground station as shown in Figure 2. The ground station sends control commands to and receive navigation data from the aircraft. In general, there exists two-way communication delay, which may cause instability to the aircraft. To guarantee aircraft safety, it is also desirable to constrain the roll angle and the pitch angle of the aircraft to be within certain ranges. In this case, the aircraft can be equipped with an onboard satisficing controller, which defines the proper satisficing set within which the aircraft maintains its base-line performance. Both the human input and the aircraft states obtained from the onboard sensor are sent to the satisficing controller. If the aircraft is safe, that is, its states are within the satisficing set, the human input will be used to control the aircraft. Otherwise, the human input will be overridden by a control element, detailed in the sequel, in the satisficing set.

Define the continuous scalar function $k : \mathbb{R}^n \to [0, 1]$ as follows:

$$k(x) = \begin{cases} 1, & x \in R_s(c - \epsilon) \\ l(x), & x \in R_s(c) \setminus R_s(c - \epsilon) \text{ and } x \notin B(R_s(c)) \\ 0, & x \notin R_s(c) \text{ or } x \in B(R_s(c)) \end{cases}$$
(8)

where $R_s(\cdot)$ and c are defined in Definition 3.2, $B(\cdot)$ represents the boundary of a set, $\epsilon = [\epsilon_1, \cdots, \epsilon_J, -\epsilon_{J+1}, \cdots, -\epsilon_K]^T$ with $0 < \epsilon_i < c_i$, $i \in \mathcal{I}$ and $\epsilon_i > 0$, $i \in \overline{\mathcal{I}}$, and $l(x) \in [0, 1]$ can be any continuous function in x. The introduction of function l(x)defines a smooth transition from human input to snapped control input when the system starts moving out of the safe region.

Note that l(x) in Eq. (8) may be defined in various ways. One option of defining l(x) is based on the following two sets of functions. Define two sets of continuous scalar functions \mathcal{F}_a and \mathcal{F}_r as follows:

$$\mathcal{F}_{a}(\epsilon_{a}, c_{a}) = \{ f : \mathbb{R}^{+} \to [0, 1] | f(r) \\ = \begin{cases} 1, & 0 \le r \le c_{a} - \epsilon_{a} \\ p(r), & c_{a} - \epsilon_{a} < r < c_{a} \} \\ 0, & r \ge c_{a} \end{cases}$$
(9)

$$\mathcal{F}_{r}(\epsilon_{r}, c_{r}) = \{ f : I\!\!R^{+} \to [0, 1] | f(r) \\ = \begin{cases} 0, & 0 \le r \le c_{r} \\ q(r), & c_{r} < r < c_{r} + \epsilon_{r} \\ 1, & r \ge c_{r} + \epsilon_{r} \end{cases}$$
(10)

where c_a and c_r are positive constants, $0 < \epsilon_a < c_a, \epsilon_r > 0$, and $p(r), q(r) \in [0, 1]$ are monotonically continuous with $p(c_a - \epsilon_a) = 1$, $p(c_a) = 0$, $q(c_r) = 0$, and $q(c_r + \epsilon_r) = 1$. Accordingly, l(x) can be defined as $l(x) = \prod_{i=1}^{K} k_{hi}(V_i(x))$, where $V_i(x)$ is the *i*th CLF given in Definition 3.2, $k_{hi} \in \mathcal{F}_a(\epsilon_i, c_i)$ for $i \in \mathcal{I}$, and $k_{hi} \in \mathcal{F}_r(\epsilon_i, c_i)$ for $i \in \overline{\mathcal{I}}$. Obviously p(r) and q(r)are easy to define in this case, for example, a linear strictly decreasing function for p(r) and a linear increasing function for q(r).

Define the control input to system (2) as

$$u = (1 - k(x))d(x) + k(x)h(t),$$
(11)

where k(x) is given by Eq. (8), h(t) is the human input, and d(x) is the satisficing control snapped from the human input to the satisficing input set U given by

$$d(x) = \arg \min_{u \in U} ||h - u||.$$
 (12)

With the asymptotically satisficing input set U_a , d(x) can also be the asymptotically satisficing control given by

$$d(x) = \arg \min_{u \in U_a} \|h - u\|.$$
 (13)

From Eq. (11), we can see that the control input to the system is the human input when the system state stays in the safe region and is switched to the satisficing control when the system state leaves the safe region. When the human input is overridden in unsafe regions, the control input u to the system approximates the human input as closely as possible from the satisficing input set. The function k(x) in Eq. (8) can be constructed based on design requirements. The tolerance vector ϵ in Eq. (8) guarantees transition from human input to the satisficing control starts when the system state is close to the boundary of the safe region.

Definition 3.4: We say that the safeguarded shared control problem is solved if the following two conditions are satisfied.

- If the system state is initially within the safe region, then the state will stay within the region for all the time and for any human input.
- If the system state is initially outside the safe region, the state will be driven into the safe region within finite time for any human input.

We have the following theorem for safeguarded shared control for static safe regions.

Theorem 3.5: If the human input h(t) is continuous (piecewise continuous), the control input to system (2) is given by Eq. (11), and d(x) is given by Eq. (13), the shared control problem is solved. Also, if d(x) is given by Eq. (12), the first condition in Definition 3.4 is satisfied.

Proof: If h(t) is continuous (piecewise continuous), we know that d(x) is also continuous (piecewise continuous) following a similar argument for the continuity of the minnorm control laws in [9]. Therefore, u is also continuous (piecewise continuous) since k(x) is always continuous. **Condition 1.**

Assume that the system state x is initially within R_s and reaches the boundary of R_s at time $t = t_1$.

From Eq. (8), we know that $k(x(t_1)) = 0$ for $x(t_1) \in B(R_s)$. Thus $u = d(x(t_1))$ at $t = t_1$ from Eq. (11).

Given Eq. (12), it is easy to see that $u \in U$ at $t = t_1$. Given Eq. (13), it is easy to see that $u \in U_a$ at $t = t_1$. From the property of the satisficing input set and asymptotically satisficing input set, we know that $\dot{V}_i(x) \leq 0$ for all $x \in$ $B_i(R_i(c_i)), i \in \mathcal{I}$, and $\dot{V}_i(x) \geq 0$ for all $x \in B_i(R_i(c_i))$, $i \in \overline{\mathcal{I}}$, where $B_i(\cdot)$ represents the boundary of the *i*th safe subregion, at $t = t_1$.

Therefore, the state x cannot leave the safe region R_s for $t \ge t_1$.

Condition 2.

Assume that the system state x is outside the safe region R_s at $t = t_0$. Similar to condition 1, we know that u = d(x) from Eqs. (8) and (11).

Given Eq. (13), it is easy to see that $u \in U_a, \forall x \notin R_s$, which implies that $\dot{V}_i(x) < 0$ for all $x \notin R_i(c_i), i \in \mathcal{I}$, and $\dot{V}_i(x) > 0$ for all $x \notin R_i(c_i), i \in \overline{\mathcal{I}}$.

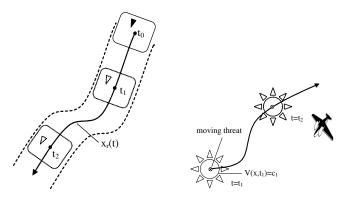
Therefore, from the standard Lyapunov theory, the state x will be driven to R_s within finite time.

B. Safeguarded Control for Dynamic Safe Region

Definition 3.6: A dynamic safe region corresponds to a time-varying region of state space where the system has guaranteed minimum base-line performance.

Similar to a static safe region, a dynamic safe region for attractive behaviors is defined as $R = \{x : V(x,t) \le c\}$, where V(x,t) = 0 at $x = x_r(t)$ and $x_r(t)$ is the reference state trajectory to follow. A dynamic safe region for repulsive behaviors is defined accordingly as $R = \{x : V(x,t) \ge c\}$, where V(x,t) = 0 at $x = x_r(t)$ and $x_r(t)$ is the reference state trajectory to avoid.

Two illustrative examples are shown in Figure 3 and 4 respectively. In Figure 3, a human controls a vehicle to follow a given trajectory $x_r(t)$. As long as the tracking error is below a certain upper bound, the human has the full authority and flexibility to maneuver the vehicle. Once the tracking error is above the upper bound, controls from the satisficing input set take over the human input and guarantee minimum tracking performance. We can see that the vehicle (denoted by a wedge) is required to stay within the rectangle area centered at the current trajectory $x_r(t_*)$ at each time t_* . In Figure 4, a human maneuvers an airplane but needs to avoid a moving threat with trajectory $x_r(t)$. Here we assume that the vehicle can measure the current position and velocity of the moving threat. We can see that the airplane



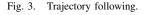


Fig. 4. Dynamic threat avoidance.

is required to stay outside the circle area centered at the current trajectory $x_r(t_*)$ at each time t_* .

For the case of dynamic safe regions, safeguarded control u is defined the same as equation (11) except that $V_i(x)$ is replaced with $V_i(x,t)$. We will not repeat the similar procedures as in Section III-A but consider safeguarded control for the special case of the trajectory following problems.

Given a time-invariant system, a common way to tackle a tracking problem is to perform a state substitution, i.e. define $\tilde{x}(t) = x(t) - x_r(t)$, where $x_r(t)$ is the reference state trajectory to follow. As a result, we have a time-varying system for \tilde{x} . Then we need to find a CLF $V(\tilde{x},t)$ for this time-varying system in order to apply the safeguarded control, which may be hard to find. Also, even if such a CLF can be found, a safeguarded control may not be implemented effectively due to the possible input constraints of the vehicle. That is, the trajectory may evolve too fast for the vehicle to track even with the maximum control effort the vehicle can have. Therefore, we introduce the following pointwise CLF and parameterized trajectory ideas motivated by [10], [11] to add feedback from the vehicle to the reference trajectory.

Definition 3.7: A continuously differentiable function $V : \mathbb{I}\!\!R^n \times \mathbb{I}\!\!R \to \mathbb{I}\!\!R$ is a pointwise control Lyapunov function (pCLF) for system (2) if it is positive definite, radially unbounded, and satisfies

$$\inf_{u} \left\{ \frac{\partial V(x,s)}{\partial x} \left(f(x,s) + g(x,s)u \right) \right\} < 0, \quad (14)$$

 $\forall x \neq 0$ and each constant $s \in [s_1, s_2]$.

Hereafter we assume that a pointwise CLF can be found for a smooth parameterized reference trajectory $x_r(s)$, that is, V(x,s) = 0 at $x = x_r(s)$ for each $s \in [s_1, s_2]$. We also define the safe region for a parameterized trajectory following as

$$R_s(c,s) = \{x : V(x,s) \le c\}, \quad s \in [s_1, s_2].$$
(15)

Unlike the static safe region in Section III-A, this safe region is time-varying with respect to the parameter s.

Define $\underline{b}(x,s)$ as

$$\underline{b}(x,s) \stackrel{\triangle}{=} \begin{cases} -\sigma(V(x,s)) / \left(\frac{\partial V(x,s)}{\partial x}f\right), & \frac{\partial V(x,s)}{\partial x}g = 0\\ \mu, & \frac{\partial V(x,s)}{\partial x}g \neq 0 \end{cases}$$

where $\sigma(\cdot)$ is a class ${\cal K}$ function and μ can be any given positive number.

The pointwise stabilizing set in this case is defined as

$$S_b(x,s) = \left\{ u : \frac{\partial V(x,s)}{\partial x} (f + gu) \le -\frac{1}{b} \sigma(V(x,s)), b \ge \underline{b} \right\},$$
(16)

where $s \in [s_1, s_2]$.

Lemma 3.1: For each $s \in [s_1, s_2]$, $\underline{b}(x)$ is positive and $S_b(x,s)$ is nonempty for each $x \neq x_r(s)$.

Proof: We follow a similar argument used in [3].

If $\frac{\partial V(x,s)}{\partial x}g = 0$, we know that $\frac{\partial V(x,s)}{\partial x}f < 0$ for each $x \neq x_r(s)$ since V(x,s) is a pointwise CLF. Thus $\underline{b}(x) = -\sigma(V(x,s))/\left(\frac{\partial V(x,s)}{\partial x}f\right) > 0$ for each $x \neq x_r(s)$.

Also, we know that

$$b \ge \underline{b}(x) > 0$$
$$\implies -\frac{1}{b}\sigma(V(x,s)) \ge -\frac{1}{\underline{b}(x)}\sigma(V(x,s)) = \frac{\partial V(x,s)}{\partial x}f$$

which means that $S_b(x, s)$ is nonempty when $\frac{\partial V(x,s)}{\partial x}g = 0$. If $\frac{\partial V(x,s)}{\partial x}g \neq 0$, we know that $\underline{b}(x) = \mu > 0$ and a control u can always be found to satisfy the inequality represented in $S_b(x, s)$.

Following [10], let \dot{s} be given by

$$\dot{s} = \begin{cases} \min\left\{\frac{v_0}{\delta + \|\frac{\partial x_r(s)}{\partial s}\|}, \frac{-(\frac{\partial V}{\partial x})^T \dot{x}}{\delta + |\frac{\partial V}{\partial s}|} \left(\frac{\sigma(c)}{\sigma(V(x,s))}\right)\right\},\\ s_1 \le s < s_2\\ 0, \qquad s = s_2 \end{cases},$$
(17)

where c is the upper bound for the pointwise CLF, $\delta > 0$ is a small positive constant, and v_0 is the nominal velocity for the reference trajectory.

We need the following lemma to show our main theorem. Lemma 3.2: If \dot{s} is given by (17), then $V(x(t_0), s(t_0)) \leq$ c implies that $V(x(t), s(t)) \leq c, \forall t \geq t_0$. Furthermore, if $s \in [s_1, s_2]$ and there is a control u(x, s) and a constant L > 0 such that $\frac{-\frac{\partial V}{\partial x} \dot{x}}{\sigma(V(x,s))} \ge L$, s will reach s_2 within finite time and $\left\|\frac{dx_r}{dt}\right\| \approx v_0$ for $V(x,s) \ll c$ and $\delta v_0 \ll \left\|\frac{\partial x_r(s)}{\partial s}\right\|$. Proof: see [10].

Lemma 3.3: Given Eq. (17), for each $s \in [s_1, s_2]$ and $x \neq x_r(s)$, controls chosen from $S_b(x,s)$ guarantee that there exists bounded $b \geq \underline{b}$ such that $\frac{-\frac{\partial V}{\partial x}x}{\sigma(V(x,s))}$ has a positive lower bound. Furthermore, controls chosen from $S_b(x,s)$ for bounded b guarantee that $V(x(t), s(t)) \leq c$ for arbitrary

 $V(x(t_0), s(t_0) > c \text{ within finite time.}$ $Proof: \text{ If } \frac{\partial V(x,s)}{\partial x}g = 0, \text{ we know that } \underline{b}(x,s) = -\sigma(V(x,s))/(\frac{\partial V(x,s)}{\partial x}f) \text{ is continuous for each } x \neq x_r(s)$ from the property of the pCLF. From the first argument in Lemma 3.2, ||x|| is bounded given Eq. (17) for each $x \neq x_r(s)$. When $x \to x_r(s)$, function $\sigma(\cdot)$ can be chosen so that $\sigma(V(x,s))$ has a higher order than $\frac{\partial V(x,s)}{\partial x}f$, that is $\begin{array}{l} \lim_{x\to x_r(s)} -\sigma(V(x,s))/(\frac{\partial V(x,s)}{\partial x}f) = 0. \mbox{ Thus } \underline{b}(x,s) \mbox{ is bounded for each } x\neq x_r(s). \\ \mbox{ If } \frac{\partial V(x,s)}{\partial x}g \eq 0, \enskip \underline{b}(x,s) \eq \mu, \mbox{ which is obviously bounded.} \end{array}$

Controls chosen from $S_b(x,s)$ satisfies $\frac{\partial V}{\partial x}\dot{x} \leq -\frac{1}{b}\sigma(V(x,s)) \iff \frac{-\frac{\partial V}{\partial x}^T\dot{x}}{\sigma(V(x,s))} \geq \frac{1}{b}$, which has a positive lower bound since an upper bounded b can be chosen to be above bounded b(x, s).

The last part of the lemma follows a similar argument for Theorem 4.4 in [12].

To choose stabilizing controls from (16), it is desirable to have an explicit expression for a locally Lipschitz continuous b.

Define b as

$$b(x,s) \stackrel{\triangle}{=} \begin{cases} -\sigma(V)/\frac{\partial V}{\partial x}f, & \frac{\partial V}{\partial x}g = 0\\ \frac{2\frac{\partial V}{\partial x}f + 2\sqrt{(\frac{\partial V}{\partial x}f)^2 + \sigma(V)|\frac{\partial V}{\partial x}g|}}{|\frac{\partial V}{\partial x}g|}, & \frac{\partial V}{\partial x}g \neq 0 \end{cases},$$
(18)

which is locally Lipschitz for each $x \neq x_r(s)$ following a similar argument in [13] and [3].

For safeguarded shared control for parameterized trajectory following, we define the control input as

$$u(x,s) = (1 - k(x,s))d(x,s) + k(x,s)h(t),$$
(19)

where h(t) is the human input, k(x,s) is defined similar to Eq. (8) with $R_s(\cdot)$ given by Eq. (15), and d(x,s) is the satisficing control snapping the human input to the satisficing input set $S_b(x,s)$ with bounded b given by $d(x,s) = \arg \min_{u \in S_h(x,s)} ||h - u||.$

We have the following theorem for safeguarded shared control for parameterized trajectory following.

Theorem 3.8: If $s \in [s_1, s_2]$, \dot{s} is given by Eq. (17), u(x, s) is given by Eq. (19), the safeguarded shared control problem for parameterized trajectory following is solved. Furthermore, the last two arguments in Lemma 3.2 are also satisfied, that is, finite completion time and given reference trajectory velocity are guaranteed.

Proof: Condition 1 for shared control problem is guaranteed by the definition of \dot{s} following Lemma 3.2. Condition 2 for shared control problem is guaranteed by the last part of Lemma 3.3. The proof for the last part of this theorem follows a combination of Lemma 3.3 and 3.2.

C. Discussion

One limitation of our approach is that CLFs may not be easy to find for general affine nonlinear systems although CLFs can be constructed for systems with certain structures [14], for example, feedback linearizable systems and systems with a cascade structure. This limitation can be somewhat mitigated by using weak CLFs if proper for some applications since weak CLFs are relatively easy to find. A concern with any approach for shared control is the computation cost. Although the safe region defined in our approach is an intersection of a group of subregions, we only need to consider the relevant subregions where the system state is close to its boundaries, which can significantly drop the computation cost.

The main theme of our approach is to preserve human intention as much as possible while guaranteeing system stability. Even with the system state outside the safe region, if the human input is within the satisficing set, it will be the input to the system. If the human input is outside the satisficing input set, the closest element in the set will be chosen as the input to the system. Therefore, in the case of human neglect, minimum performance (e.g. system stability) is still guaranteed.

IV. CONCLUSION

This paper has presented a preliminary result for safeguarded shared control using rigorous Lyapunov stability perspective. We applied control Lyapunov functions to define static and dynamic safe regions and also satisficing input sets. The control input was defined as a combination of the human input and the satisficing control element that is closest to the human input in the satisficing set. We showed that this approach not only gives humans maximum flexibility but guarantees system stability and minimum base-line performance. The readers can refer to [15] for a simulation study where a mobile robot effectively achieved wall avoidance and obstacle avoidance.

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